

Completely solving general SDPs

Bruno F. Lourenço
ISM

November 27th, 2021
T60

Joint work with Masakazu Muramatsu and Takashi Tsuchiya

SDPs

$$\begin{array}{ll}
 \inf_x \langle c, x \rangle & \text{(P)} \\
 \text{subject to } \mathcal{A}x = b \\
 x \in \mathcal{S}_+^n
 \end{array}
 \qquad
 \begin{array}{ll}
 \sup_y \langle b, y \rangle & \text{(D)} \\
 \text{subject to } c - \mathcal{A}^*y \in \mathcal{S}_+^n,
 \end{array}$$

How to solve SDPs in *general*?



B. F. Lourenço, M. Muramatsu, and T. Tsuchiya,

Solving SDP completely with an interior point oracle

Optimization Methods and Software, 36 (2021), pp. 425–471.

Strange behaviour 1 - Duality gaps

$$\begin{array}{ll}
 \inf_x & \begin{array}{l} \cancel{-2x_{12}} \xrightarrow{0} \cancel{-2x_{13}} \xrightarrow{0} \\ \text{s.t. } x_{11} = 0 \\ \quad -x_{22} - 2x_{13} = -1 \\ \quad x \in \mathcal{S}_+^3. \end{array} & \text{(P)}
 \end{array}
 \qquad
 \begin{array}{ll}
 \sup_{t,s} & \begin{array}{l} \xrightarrow{s} \xrightarrow{-1} \\ \text{s.t. } \begin{pmatrix} t & 1 & s-1 \\ 1 & s & 0 \\ s-1 & 0 & 0 \end{pmatrix} \in \mathcal{S}_+^3 \end{array} & \text{(D)}
 \end{array}$$

$$\theta_D = -1 \text{ and } \theta_P = 0.$$

Strange behaviour 2 - Non-attainment

$$\begin{aligned} \sup_{t,s} \quad & -s && \text{(D)} \\ \text{s.t.} \quad & \begin{pmatrix} t & 1 \\ 1 & s \end{pmatrix} \in \mathcal{S}_+^2 \end{aligned}$$

Strange behaviour 3 - Weak infeasibility

$$\begin{aligned} \sup_{t,s} \quad & t && \text{(D)} \\ \text{s.t.} \quad & \begin{pmatrix} t & 1 \\ 1 & 0 \end{pmatrix} \in \mathcal{S}_+^2 \end{aligned}$$

- Let $V = \{c - \mathcal{A}^*y \mid y \in \mathbb{R}^n\}$
- In general, (D) feasible $\Rightarrow \text{dist}(V, \mathcal{S}_+^n) = 0$
- Here, we have $\text{dist}(V, \mathcal{S}_+^n) = 0$, but (D) is **infeasible**.

How to solve SDPs in *general*?

- ① IPMs? Some first order method? **Probably won't work if there is positive duality gap or non-attainment**
- ② What if we try to regularize the SDP via *facial reduction* or something?
 - ① It only fixes **one side of the problem**.

It is very hard to solve general SDPs! Even in low-dimensions and with apparently harmless data...

Ok, so which SDPs can we actually solve?

- If (P) and (D) **both** have interior points, then $\theta_P = \theta_D$ and are attained.
 - We have a decent chance of actually solving (P) and (D) with IPMs, augmented Lagrangian and etc.

The interior point oracle \mathcal{O}_{int}

Input: The problem data: \mathcal{A}, b, c . Both (P) and (D) must have interior points.

Output: A primal-dual optimal solution pair x^*, y^* .

The main result

Completely solving SDPs

Any SDP can be **completely solved** via polynomially (in n) many calls to \mathcal{O}_{int}

Completely solving (D) entails the following.

- Deciding feasibility and infeasibility.
 - In case of infeasibility, distinguishing between weak and strong infeasibility.
- Computing the optimal value
 - If attained, we also want an optimal solution.
 - If not, we compute an ϵ -optimal solution for any $\epsilon > 0$.
 - We also want to detect unboundedness.

Next we describe our tools: facial reduction and double facial reduction.

Facial Reduction Basics

$$\begin{aligned} & \sup_y \langle b, y \rangle && (D) \\ & \text{subject to } c - \mathcal{A}^*y \in \mathcal{S}_+^n, \end{aligned}$$

Let \mathcal{F}_D denote the feasible slacks of (D), $\mathcal{F}_D = \{S \in \mathcal{S}_+^n \mid \exists y, c - \mathcal{A}^*y\}$

- If \mathcal{F}_D has no interior point of \mathcal{S}_+^n then \mathcal{F}_D lies on a proper face of \mathcal{S}_+^n
- The smallest such face $\mathcal{F} \trianglelefteq \mathcal{S}_+^n$ contains \mathcal{F}_D and

$$\mathcal{F}_D \cap \text{ri } \mathcal{F} \neq \emptyset.$$

- Replacing \mathcal{S}_+^n by \mathcal{F} leads to a **smaller equivalent** problem that has an interior point!



J. M. Borwein and H. Wolkowicz.

Regularizing the abstract convex program.

Journal of Mathematical Analysis and Applications, 83(2):495 – 530, 1981.

More about facial reduction

- If $(c + \text{range } \mathcal{A}^*) \cap \text{ri } \mathcal{S}_+^n = \emptyset$, we find a hyperplane $\{d\}^\perp$ that properly separates both, with $d \in \mathcal{S}_+^n$.
- Then, we replace \mathcal{S}_+^n by $\mathcal{S}_+^n \cap \{d\}^\perp$ and repeat.

Example:

$$\begin{pmatrix} t & 1 & s-1 \\ 1 & s & 0 \\ s-1 & 0 & 0 \end{pmatrix} \in \mathcal{S}_+^3$$

We can let $d = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$

Facial reduction and \mathcal{O}_{int}

Theorem

Through $\mathcal{O}(n)$ calls to \mathcal{O}_{int} we can either detect that (D) is infeasible or find an equivalent SDP that has an interior point at the dual side.

Key idea: d can be found by solving by successively using \mathcal{O}_{int} to solve

$$\begin{aligned} & \inf_{x, t, w} t && (P_{\mathcal{K}}) \\ \text{subject to} & -\langle c, x - te^* \rangle + t - w && = 0 && (1) \\ & \langle e, x \rangle + w && = 1 && (2) \\ & \mathcal{A}x - t\mathcal{A}e^* && = 0 && (3) \\ & (x, t, w) \in \mathcal{K}^* \times \mathbb{R}_+ \times \mathbb{R}_+ \end{aligned}$$

$$\begin{aligned} & \sup_{y_1, y_2, y_3} y_2 && (D_{\mathcal{K}}) \\ \text{subject to} & cy_1 - ey_2 - \mathcal{A}^*y_3 \in \mathcal{K} && (4) \\ & 1 - y_1(1 + \langle c, e^* \rangle) + \langle e^*, \mathcal{A}^*y_3 \rangle \geq 0 && (5) \\ & y_1 - y_2 \geq 0 && (6) \end{aligned}$$

with $\mathcal{K} = \mathcal{S}_+^n$, $\mathcal{K} = \mathcal{S}_+^n \cap \{d_1\}^\perp$, $\mathcal{K} = \mathcal{S}_+^n \cap \{d_2\}^\perp$ and so on.

The story so far

Suppose we wish to solve (D)

$$\begin{aligned} & \sup_y \langle b, y \rangle && \text{(D)} \\ & \text{subject to } c - \mathcal{A}^*y \in \mathcal{S}_+^n, \end{aligned}$$

From facial reduction we either detect infeasibility or obtain

$$\begin{aligned} & \sup_y \langle b, y \rangle && (\hat{D}) && \inf_x \langle c, x \rangle && (\hat{P}) \\ & \text{subject to } c - \mathcal{A}^*y \in \mathcal{F}_{\min}^D. && && \text{subject to } \mathcal{A}x = b \\ & && && x \in (\mathcal{F}_{\min}^D)^* \end{aligned}$$

where $\mathcal{F}_{\min}^D \subseteq \mathcal{S}_+^n \subseteq (\mathcal{F}_{\min}^D)^*$

(\hat{D}) is equivalent to (D) and has (relative) interior points.

However we can not use \mathcal{O}_{int} to solve (\hat{P}) and (\hat{D}) yet.

Double facial reduction

Idea: apply facial reduction to (\hat{P}) .

$$\begin{aligned} & \sup_y \langle b, y \rangle && (D) \\ & \text{subject to } c - \mathcal{A}^* y \in \mathcal{S}_+^n, \end{aligned}$$

First FR

$$\begin{array}{ll} \sup_y \langle b, y \rangle & (\hat{D}) \\ \text{subject to } c - \mathcal{A}^* y \in \mathcal{F}_{\min}^D. & \end{array} \qquad \begin{array}{ll} \inf_x \langle c, x \rangle & (\hat{P}) \\ \text{subject to } \mathcal{A}x = b & \\ x \in (\mathcal{F}_{\min}^D)^* & \end{array}$$

Second FR

$$\begin{array}{ll} \sup_y \langle b, y \rangle & (D^*) \\ \text{subject to } c - \mathcal{A}^* y \in (\mathcal{F}_{\min}^{\hat{P}})^*. & \end{array} \qquad \begin{array}{ll} \inf_x \langle c, x \rangle & (P^*) \\ \text{subject to } \mathcal{A}x = b & \\ x \in \mathcal{F}_{\min}^{\hat{P}} & \end{array}$$

The double FR theorem

$$\begin{array}{ll} \sup_y \langle b, y \rangle & (\hat{D}) \\ \text{subject to } c - \mathcal{A}^* y \in \mathcal{F}_{\min}^D. \end{array}$$

$$\begin{array}{ll} \inf_x \langle c, x \rangle & (\hat{P}) \\ \text{subject to } \mathcal{A}x = b \\ x \in (\mathcal{F}_{\min}^D)^* \end{array}$$

$$\begin{array}{ll} \sup_y \langle b, y \rangle & (D^*) \\ \text{subject to } c - \mathcal{A}^* y \in (\mathcal{F}_{\min}^{\hat{P}})^*. \end{array}$$

$$\begin{array}{ll} \inf_x \langle c, x \rangle & (P^*) \\ \text{subject to } \mathcal{A}x = b \\ x \in \mathcal{F}_{\min}^{\hat{P}} \end{array}$$

We have $\mathcal{F}_{\min}^D \subseteq (\mathcal{F}_{\min}^{\hat{P}})^*$, $\mathcal{F}_{\min}^{\hat{P}} \subseteq (\mathcal{F}_{\min}^D)^*$.

Theorem

- ① (θ_D) is finite $\iff \mathcal{F}_{\min}^{\hat{P}} \neq \emptyset$. In this case, (P^*) and (D^*) both have relative interior points and

$$\theta_D = \theta_{P^*} = \theta_{D^*}.$$

- ② $\theta_D = +\infty$ if and only if $\mathcal{F}_{\min}^{\hat{P}} = \emptyset$.

The story so far - Part 2

$$\begin{aligned} & \sup_y \langle b, y \rangle && \text{(D)} \\ & \text{subject to } c - \mathcal{A}^* y \in \mathcal{S}_+^n, \end{aligned}$$

$$\begin{aligned} & \sup_y \langle b, y \rangle && \text{(D}^*) \\ & \text{subject to } c - \mathcal{A}^* y \in (\mathcal{F}_{\min}^{\hat{P}})^*. \end{aligned} \qquad \begin{aligned} & \inf_x \langle c, x \rangle && \text{(P}^*) \\ & \text{subject to } \mathcal{A}x = b \\ & x \in \mathcal{F}_{\min}^{\hat{P}} \end{aligned}$$

So far, we are able to

- Determine whether (D) is feasible or not.
- Compute θ_D and determine whether it is $+\infty$ or not.
- **Next steps:** obtaining optimal solutions.

Optimal solutions

If we know θ_D , we can solve the feasibility problem

$$\begin{array}{ll} \text{find} & y \\ \text{subject to} & c - \mathcal{A}^*y \in \mathcal{S}_+^n, \\ & \langle b, y \rangle = \theta_D \end{array} \quad (\text{Feas})$$

in $O(n)$ calls to \mathcal{O}_{int} .

Unattained optimal solutions

$$\begin{aligned} & \sup_y \langle b, y \rangle && (\hat{D}) \\ & \text{subject to } c - \mathcal{A}^* y \in \mathcal{F}_{\min}^D. \end{aligned}$$

$$\begin{aligned} & \inf_x \langle c, x \rangle && (\hat{P}) \\ & \text{subject to } \mathcal{A}x = b \\ & && x \in (\mathcal{F}_{\min}^D)^* \end{aligned}$$

$$\begin{aligned} & \sup_y \langle b, y \rangle && (D^*) \\ & \text{subject to } c - \mathcal{A}^* y \in (\mathcal{F}_{\min}^{\hat{P}})^*. \end{aligned}$$

$$\begin{aligned} & \inf_x \langle c, x \rangle && (P^*) \\ & \text{subject to } \mathcal{A}x = b \\ & && x \in \mathcal{F}_{\min}^{\hat{P}} \end{aligned}$$

Let y^* be an optimal solution to (D^*)

The directions $\{d_1, \dots, d_\ell\}$ obtained in the final FR and y^* can be used to construct y_ϵ :

$$c - \mathcal{A}^* y_\epsilon \in \mathcal{S}_+^n, \quad \langle b, y_\epsilon \rangle \geq \theta_D - \epsilon.$$

Completely solving general SDPs

Using FR, double FR and \mathcal{O}_{int} we can

- Detect feasibility and infeasibility.
 - in case of infeasibility: detect the type of feasibility.
- Compute the optimal value and an optimal solution if it exists.
- Compute ϵ -optimal solutions.

An example

$$\sup_{y \in \mathbb{R}^8} -y_4 - 2y_6 - 2y_7 \quad (D)$$

s.t.

$$\begin{pmatrix} y_1 & & & & & & & & y_3 - 1 \\ & y_1 & & & & & & & y_5 - 1 \\ & & y_2 & y_3 & & & & & \\ & & y_3 & y_4 - y_5 & & & & & \\ & & & & y_4 & -0.5y_8 + 0.5 & y_6 & & \\ & & & & -0.5y_8 + 0.5 & y_8 & y_7 & & \\ & & & & y_6 & y_7 & 0 & & \\ y_3 - 1 & y_5 - 1 & & & & & & & 0 \end{pmatrix} \in \mathcal{S}_+^8.$$

- $\theta_D = -1$, $\theta_P = 0$ and **neither are attained.**

First FR

Let $\mathcal{S}_+^{r,n} := \left\{ \begin{pmatrix} U & 0 \\ 0 & 0 \end{pmatrix} \in \mathcal{S}^n \mid U \in \mathcal{S}_+^r \right\}$.

- $\mathcal{F}_{\min}^D = \mathcal{S}_+^{6,8}$

- $(\mathcal{F}_{\min}^D)^* = \left\{ \begin{pmatrix} U & V \\ V & W \end{pmatrix} \in \mathcal{S}^8 \mid U \in \mathcal{S}_+^6 \right\}$

Second FR

Linear constraints of (P):

$$\begin{array}{llll}
 -x_{11} - x_{22} = 0 & -x_{33} = 0 & -2x_{18} - 2x_{34} = 0 & -x_{44} - x_{55} = -1 \\
 -2x_{28} + x_{44} = 0 & -2x_{57} = -2 & -2x_{67} = -2 & x_{56} - x_{66} = 0
 \end{array}$$

x is feasible for (P) if and only if $x \in \mathcal{S}_+^8$ and can be written as

$$\begin{pmatrix}
 0_4 & & & & \\
 & 1 & x_{56} & 1 & x_{58} \\
 & x_{56} & x_{56} & 1 & x_{68} \\
 & 1 & 1 & x_{77} & x_{78} \\
 & x_{58} & x_{68} & x_{78} & x_{88}
 \end{pmatrix}$$

x is feasible for (\hat{P}) if and only if $x \in (\mathcal{S}^{6,8})^*$, satisfies the linear equations above and can be written as

$$\begin{pmatrix}
 0_3 & 0 & V_1 \\
 0 & U & V_2 \\
 V_1^T & V_2^T & W
 \end{pmatrix},$$

where $U \in \mathcal{S}_+^n$.

Second FR - (cont.)

x is feasible for $(\hat{P}) \Rightarrow x$ can be written as

$$\begin{pmatrix} 0_3 & 0 & V_1 \\ 0 & U & V_2 \\ V_1^T & V_2^T & W \end{pmatrix},$$

where $U \in \mathcal{S}_+^n$.

$$\mathcal{F}_{\min}^{\hat{P}} = \left\{ \begin{pmatrix} 0_3 & 0 & V_1 \\ 0 & U & V_2 \\ V_1^T & V_2^T & W \end{pmatrix} \mid U \in \mathcal{S}_+^3 \right\}$$
$$(\mathcal{F}_{\min}^{\hat{P}})^* = \left\{ \begin{pmatrix} Z_1 & Z_{12} & 0 \\ Z_{12}^T & U & 0 \\ 0 & 0 & 0 \end{pmatrix} \mid U \in \mathcal{S}_+^3 \right\}$$

Second FR - (cont.)

$$\begin{aligned} \sup_{y \in \mathbb{R}^8} \quad & -y_4 - 2y_6 - 2y_7 & (D^*) \\ \text{s.t.} \quad & \end{aligned}$$

$$\left(\begin{array}{cccccccc} y_1 & & & & & & & y_3 - 1 \\ & y_1 & & & & & & y_5 - 1 \\ & & y_2 & y_3 & & & & \\ & & y_3 & y_4 - y_5 & & & & \\ & & & & y_4 & -0.5y_8 + 0.5 & y_6 & \\ & & & & -0.5y_8 + 0.5 & y_8 & y_7 & \\ & & & & y_6 & y_7 & 0 & \\ y_3 - 1 & y_5 - 1 & & & & & & 0 \end{array} \right) \in (\mathcal{F}_{\min}^{\hat{P}})^*.$$

We have $y_3 = y_5 = 1$, $y_7 = y_6 = 0$. We can take $y_4 = y_5 = 1$, $y_8 = 1$ and $y_1 = y_2 = 0$. The optimal value is attained, but is not feasible for (D).

Almost optimal solution

Let $y_1^1 = y_2^1 = 1$ and $y_3^1 = \dots = y_8^1 = 0$ and let f^1 the corresponding matrix in range \mathcal{A} , so that $f_1 = \begin{pmatrix} l_3 & 0 \\ 0 & 0_5 \end{pmatrix}$

Let

$$\hat{y} = (\hat{y}_1, \hat{y}_2, \hat{y}_3, \hat{y}_4, \hat{y}_5, \hat{y}_6, \hat{y}_7, \hat{y}_8) = (1, 2, 1, 2, 1, 0, 0, 1),$$

so that $c - A^* \hat{y} \in \text{ri } \mathcal{F}_{\min}^D$.

Suppose $\epsilon = 0.1$. Let

$$\beta = \frac{\theta_D - \langle b, \hat{y} \rangle - \epsilon}{\theta_D - \langle b, \hat{y} \rangle} = 0.9.$$

Then,

$$\tilde{y} = \beta y_i^* + (1 - \beta) \hat{y} + 10y^1$$

is an ϵ -optimal solution to (D).

Conclusion

- 1 A general SDP **can** be completely solved if you are only allowed to solve SDPs having interior points at the primal and dual sides.
- 2 More stuff in the paper! Ex: discussion on different types of infeasibility, in-depth analysis of double facial reduction and more.
- 3 The results are valid for **any closed convex cone**.



B. F. Lourenço, M. Muramatsu, and T. Tsuchiya,
Solving SDP completely with an interior point oracle
Optimization Methods and Software, 36 (2021), pp. 425–471.